

Assignment 4

Post Date: 12 Nov 2018 **Due Date:** 19 Nov 2018, 11:30 am
You are permitted and encouraged to work in groups of two.

Problem 1: Union-Find with Path Compression

5 Points

Consider FIND with the following alternative path compression: After traversing the path from a vertex to its root, we update the parent pointer of each vertex along the path to point to its grandparent. Consider, e.g., subpath

$$i \rightarrow j \rightarrow k \rightarrow l \rightarrow \dots$$

Performing FIND(i) with alternative path compression results in k being predecessor of i and l being predecessor of j . Direct successors of the root keep the root as predecessor.

Go through the proof of the *Theorem of Hopcroft & Ullman* and find the inferences that require FIND to be implemented with path compression. Is the proof still correct if the alternative path compression is used?

Problem 2: Independent Vertex Sets

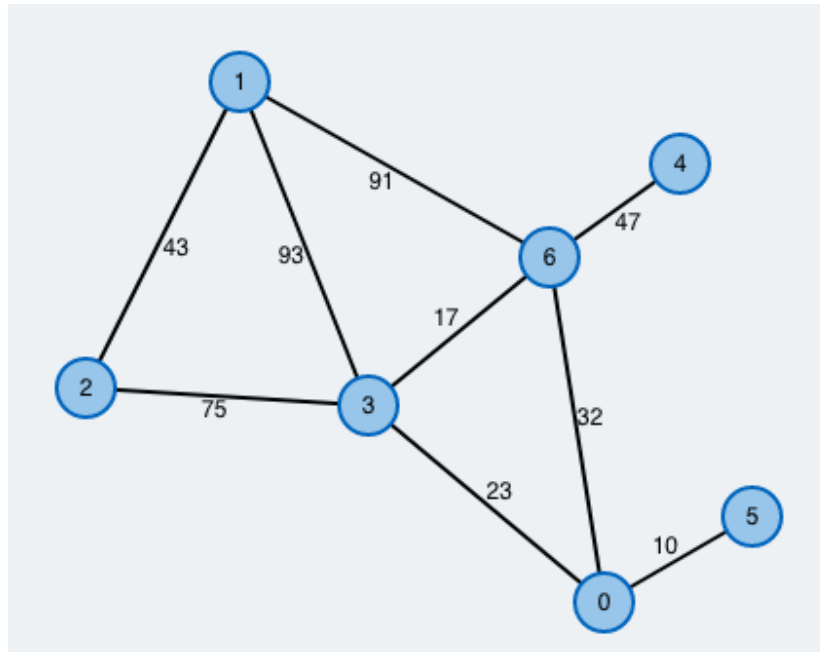
5 Points

Let $G = (V, E)$ be a graph. Let $\mathcal{I} = \{V' \subseteq V; \{u, v\} \notin E \text{ for } u, v \in V'\}$.

- (a) Show that (V, \mathcal{I}) is an independence system.
- (b) Is (V, \mathcal{I}) also a matroid? Justify your answer.

Problem 3: Minimum Spanning Tree**4 Points**

Create manually the minimum spanning tree of the graph below. Explain your algorithm and why it is an application of red and blue rules.

**Problem 4: Unique Minimum Spanning Trees****6 Points**

- Prove: Let G be a connected graph with real-valued edge weights. If for each cut the crossing edge with the lightest weight is unique, then G has a unique minimum spanning tree.
- Does the inverse of the implication in (a) hold?
- Prove or disprove: If in a connected graph G with real-valued edge weights all edges have pairwise distinct edge weights, then G has a unique minimum spanning tree.
- Does the inverse of the implication in (c) hold?